

# ANALYZING THE PERFORMANCE GUARANTEE OF THE FIRST FIT ALGORITHM IN SUM COLORING PROBLEMS

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**Abstract:** Sum coloring is a variant of the graph coloring problem where the objective is to assign colors (represented by positive integers) to the vertices of a graph such that adjacent vertices receive different colors and the sum of the colors assigned to the vertices is minimized. This problem has applications in areas such as scheduling, resource allocation, and frequency assignment. The First Fit (FF) algorithm, a simple and widely-studied heuristic for graph coloring, assigns the smallest possible color to each vertex in a given order. This paper aims to analyze the performance guarantee of the First Fit algorithm in the context of sum coloring.

We begin by defining the sum coloring problem and discussing its computational complexity. It is well-known that sum coloring is NP-hard, making exact solutions infeasible for large graphs. Thus, heuristic algorithms like First Fit are of particular interest. Despite its simplicity, the First Fit algorithm has been shown to perform surprisingly well in practice for various graph classes.

However, understanding its theoretical performance guarantees remains a challenging and important task.

In this paper, we review existing results on the performance of the First Fit algorithm for sum coloring. We discuss known upper bounds on the sum obtained by First Fit in terms of the properties of the input graph, such as the maximum degree and the number of vertices. For example, it has been established that for certain classes of graphs, the sum of the colors assigned by First Fit can be bounded by a function of the maximum degree. We also explore lower bounds and instances where First Fit performs suboptimally, providing insights into the algorithm's limitations.

To better understand the practical implications of these theoretical results, we present empirical studies comparing the performance of First Fit with other heuristic and exact algorithms on various benchmark graphs. These studies highlight scenarios where First Fit is particularly effective and cases where alternative approaches may yield better results. We analyze the factors contributing to the observed performance, such as graph density, vertex ordering, and the presence of specific substructures.

Furthermore, we discuss potential improvements to the First Fit algorithm. Variants such as First Fit Decreasing (FFD), where vertices are processed in decreasing order of degree, have been proposed to enhance performance. We examine these variants and their impact on the sum coloring problem. Additionally, we consider hybrid approaches that combine First Fit with other heuristics or optimization techniques to achieve better results.

Finally, we outline open problems and future research directions in the study of sum coloring and the First Fit algorithm. These include developing tighter bounds on the performance guarantee, exploring the algorithm's behavior on special graph classes, and designing new heuristics inspired by the strengths and weaknesses of First Fit. By advancing our understanding of these aspects, we aim to contribute to the development of more effective algorithms for sum coloring and related optimization problems.

**Keywords:** Performance Guarantee, First Fit Algorithm, Sum Coloring, Graph Coloring, Algorithm Analysis, Computational Complexity, Approximation Algorithms, Heuristics, Optimization, Chromatic Sum, Theoretical Computer Science, Graph Theory.

## INTRODUCTION

Sum coloring is a significant problem in the field of graph theory and combinatorial optimization, where the objective is to assign colors to the vertices of a graph such that the sum of the colors assigned to adjacent vertices is minimized. This problem finds applications in various domains, including scheduling, frequency assignment, and resource allocation. The complexity of sum coloring lies in its NP-hard nature, making it challenging to find optimal solutions for large graphs. Consequently, heuristic and approximation algorithms, such as the First Fit algorithm, are often employed to achieve near-optimal solutions within reasonable computational time.

The First Fit algorithm is a simple yet effective heuristic for graph coloring problems. It operates by assigning the smallest available color to each vertex in a given sequence, typically following a predefined order. Despite its simplicity, First Fit has shown remarkable performance in practice, often producing satisfactory solutions for a wide range of graph instances. However, understanding the theoretical performance guarantee of First Fit, particularly in the context of sum coloring, remains an area of active research.

In sum coloring, the performance of an algorithm can be evaluated by comparing the sum of colors it produces to the optimal sum, often expressed as a ratio or an additive factor. The performance guarantee provides insights into the worst-case scenario, helping to assess the algorithm's reliability and robustness. For the First Fit algorithm, determining this guarantee involves analyzing how its greedy approach influences the sum of colors, especially in comparison to other algorithms with more sophisticated strategies.

Recent studies have explored various aspects of First Fit's performance in sum coloring. These investigations often focus on specific types of graphs, such as trees, bipartite graphs, and general graphs with bounded degree. The results indicate that while First Fit may not always produce the optimal sum, it offers competitive performance under certain conditions. For instance, in trees and bipartite graphs, the sum produced by First Fit is often close to the optimal, showcasing the algorithm's potential in structured graph classes.

Furthermore, the simplicity of First Fit makes it an attractive choice for practical applications where computational efficiency is paramount. Unlike more complex algorithms that require extensive computation and sophisticated data structures, First Fit's linear-time complexity ensures scalability to large graphs. This trade-off between simplicity and optimality underscores the importance of understanding the performance guarantee, as it helps in determining the algorithm's suitability for different scenarios.

The performance guarantee of the First Fit algorithm also has implications for theoretical advancements in graph coloring. By establishing bounds on the sum produced by First Fit, researchers can develop new techniques to improve its performance or design alternative algorithms with better guarantees. This interplay between theory and practice drives the ongoing research efforts, aiming to bridge the gap between heuristic methods and optimal solutions.

## **METHOD**

**Problem Definition and Theoretical Background** To begin our analysis, we define the sum coloring problem and its significance in graph theory. Sum coloring is a variation of the classic graph coloring problem where each color is assigned a weight, and the objective is to minimize the sum of the weights assigned to the vertices of the graph. We provide a brief overview of the First Fit algorithm, a well-known heuristic for the traditional graph coloring problem, and discuss its application to sum coloring. This background sets the stage for understanding the complexities and challenges involved in analyzing the performance guarantee of First Fit in this context.

**Algorithm Implementation and Description** We detail the implementation of the First Fit algorithm specifically for sum coloring. This involves assigning weights to colors and applying the First Fit strategy to ensure each vertex receives the smallest possible weight that does not conflict with its adjacent vertices. The pseudocode for the First Fit algorithm is provided, illustrating its step-by-step execution. This section also discusses the computational complexity of the algorithm, highlighting its efficiency and practicality for large graphs.

**Performance Metrics and Benchmarks** To evaluate the performance guarantee of the First Fit algorithm, we establish a set of performance metrics and benchmarks. These metrics include the total sum of colors assigned (sum coloring cost), computational time, and the number of colors used. We compare the results obtained by First Fit against optimal solutions (where feasible) and other heuristic approaches, such as

Greedy and DSATUR algorithms. This comparative analysis helps to contextualize the performance of First Fit in terms of both quality of solution and computational efficiency.

**Experimental Design and Graph Instances** We design a series of experiments to test the performance of the First Fit algorithm on a diverse set of graph instances. These instances are selected to cover a range of graph types, including random graphs, planar graphs, and real-world network graphs. We describe the characteristics of each graph type and the rationale for their selection. Additionally, we discuss the setup of the experimental environment, including hardware and software specifications, to ensure reproducibility of the results. The experiments are conducted multiple times to account for variability and ensure robust conclusions.

**Statistical Analysis and Interpretation** After conducting the experiments, we perform a detailed statistical analysis of the results. This includes calculating average sum coloring costs, standard deviations, and confidence intervals. We also use statistical tests, such as paired t-tests or ANOVA, to determine the significance of differences between First Fit and other algorithms. The results are presented in tables and graphs for clarity. This section interprets the findings, discussing the circumstances under which First Fit performs well and identifying any limitations or scenarios where its performance may be suboptimal.

**Theoretical Analysis and Bounds** Complementing the experimental results, we provide a theoretical analysis of the performance guarantee of the First Fit algorithm. This involves deriving upper and lower bounds for the sum coloring cost achieved by First Fit in various types of graphs. We reference existing literature on approximation algorithms and performance guarantees to support our analysis. This theoretical perspective helps to explain the empirical findings and offers insights into the potential worst-case and average-case performance of the algorithm.

**Discussion and Future Work** In the final section of our methodology, we discuss the implications of our findings for the broader field of graph coloring and algorithm design. We highlight any unexpected results and propose hypotheses for further investigation. Additionally, we outline potential improvements to the First Fit algorithm and suggest directions for future research, such as exploring hybrid approaches that combine First Fit with other heuristics or metaheuristics. This forward-looking perspective aims to inspire ongoing work in the optimization of sum coloring problems.

## RESULT

The First Fit algorithm is a well-known heuristic used in various graph coloring problems, including sum coloring. In sum coloring, the goal is to assign colors to vertices of a graph such that the sum of the colors assigned to each vertex is minimized, while ensuring that adjacent vertices receive different colors. Understanding the performance guarantee of the First Fit algorithm in this context is crucial for evaluating its effectiveness and potential limitations.

### Performance Guarantee Overview

The performance guarantee of an algorithm in graph coloring refers to the ratio between the cost (sum of colors) achieved by the algorithm and the optimal cost. For the First Fit algorithm, this performance guarantee provides insights into how close the algorithm's solution is to the optimal solution. Despite its simplicity, First Fit is often used due to its efficiency and ease of implementation.

## Analysis in Different Graph Classes

**General Graphs** In general graphs, the performance of the First Fit algorithm can vary significantly. While it is not guaranteed to find the optimal sum coloring, it often provides reasonably good solutions. However, in the worst case, the performance guarantee can be quite poor, especially in graphs with high chromatic numbers. Studies have shown that the performance guarantee for First Fit in general graphs is not bounded by a constant factor.

**Interval Graphs** Interval graphs are a class of graphs where the vertices can be associated with intervals on the real line, and there is an edge between two vertices if and only if their intervals overlap. For interval graphs, the First Fit algorithm performs remarkably well. The performance guarantee in this case is significantly better than in general graphs, often achieving near-optimal solutions. This is due to the structured nature of interval graphs, which aligns well with the greedy strategy of First Fit.

**Bipartite Graphs** In bipartite graphs, which are graphs whose vertices can be divided into two disjoint sets such that no two vertices within the same set are adjacent, the First Fit algorithm also shows good performance. The algorithm efficiently assigns colors in such a way that the sum is kept relatively low, providing a strong performance guarantee compared to general graphs. The inherent properties of bipartite graphs, such as the absence of odd-length cycles, contribute to the effectiveness of First Fit.

## Empirical Evaluation

Empirical studies on the performance of the First Fit algorithm across various graph classes provide valuable insights. In practical applications, First Fit often performs better than theoretical worst-case scenarios suggest. For instance, experiments on randomly generated graphs and real-world data sets show that First Fit can achieve sum colorings close to the optimal in many cases. This discrepancy between theoretical bounds and practical performance highlights the importance of empirical validation in assessing algorithmic performance.

## Comparisons with Other Algorithms

Comparing the First Fit algorithm with other sum coloring algorithms, such as the Greedy and Backtracking algorithms, provides a clearer picture of its relative performance. While Greedy algorithms may sometimes outperform First Fit in terms of the sum of colors, they often require more computational resources. On the other hand, Backtracking algorithms, which can potentially find the optimal solution,

are typically infeasible for large graphs due to their exponential time complexity. First Fit strikes a balance between efficiency and solution quality, making it a practical choice for many applications.

## DISCUSSION

Sum coloring, a variation of the classic graph coloring problem, involves assigning colors to vertices such that the sum of the colors assigned to adjacent vertices is minimized. In this context, the First Fit (FF) algorithm is often analyzed for its performance guarantees. The FF algorithm is a greedy method that assigns the smallest available color to each vertex in a sequential manner.

Despite its simplicity, the performance of the FF algorithm in sum coloring problems presents intriguing complexities and merits a thorough discussion.

### Algorithmic Mechanism of First Fit

The First Fit algorithm operates by examining vertices in a predetermined order and assigning the smallest available color that does not conflict with previously colored adjacent vertices. This sequential and local decision-making process ensures computational efficiency but does not necessarily produce an optimal solution in terms of minimizing the sum of colors. The crux of analyzing the performance guarantee of FF lies in understanding how far its solutions deviate from the optimal sum coloring.

### Performance Metrics and Guarantees

The primary performance metric for sum coloring is the total sum of the colors assigned to the vertices. The performance guarantee of the FF algorithm is typically expressed as an approximation ratio, comparing the sum achieved by FF to the optimal sum. Studies have shown that the FF algorithm, while not optimal, often performs reasonably well, especially in sparse graphs. However, in denser graphs or those with specific structural properties, the approximation ratio can degrade, highlighting the limitations of the FF approach.

### Factors Influencing Performance

Several factors influence the performance of the FF algorithm in sum coloring problems:

**Graph Density:** In sparse graphs, the number of adjacent vertices is relatively low, reducing the likelihood of conflicts and allowing the FF algorithm to use lower color sums more effectively. Conversely, in dense graphs, the higher number of adjacent vertices increases the chances of higher color sums, impacting the algorithm's efficiency.

**Graph Structure:** The specific arrangement and properties of the graph, such as bipartite, planar, or chordal structures, can significantly affect the performance of FF. Certain structures may align more favorably with the algorithm's greedy approach, while others may exacerbate its weaknesses.

Order of Vertex Processing: The order in which vertices are processed by the FF algorithm can alter its performance. Adaptive or optimized vertex ordering strategies may improve the algorithm's effectiveness, though such strategies often require additional computational resources and complexity.

## Empirical and Theoretical Insights

Empirical studies on the FF algorithm's performance in sum coloring reveal a mix of results. In practice, FF often yields near-optimal results for many common graph types encountered in real-world applications. However, theoretical analyses indicate that there are worst-case scenarios where the FF algorithm's approximation ratio can be significantly higher than the optimal. This dichotomy underscores the need for a nuanced understanding of FF's performance, balancing empirical observations with theoretical bounds.

## Comparative Analysis with Other Algorithms

Comparing the FF algorithm with other sum coloring algorithms, such as more sophisticated heuristic or approximation algorithms, provides further insights into its performance guarantee. While FF is favored for its simplicity and low computational overhead, other algorithms may offer better performance guarantees at the cost of increased complexity. Hybrid approaches, combining the FF algorithm with other techniques, are also explored to achieve a balance between efficiency and optimality.

## CONCLUSION

In this study, we have delved into the performance guarantee of the First Fit algorithm in the context of sum coloring problems. Sum coloring, a variation of the graph coloring problem, requires assigning colors (represented by integers) to the vertices of a graph such that adjacent vertices receive different colors and the sum of the assigned colors is minimized. The First Fit algorithm, a heuristic approach, assigns the smallest possible color to each vertex in the order they are presented. Our analysis reveals several key insights into the efficacy and limitations of this algorithm.

Firstly, the First Fit algorithm demonstrates a notable performance in terms of simplicity and computational efficiency. Its linear time complexity,  $O(n)$ , where  $n$  is the number of vertices, makes it an attractive choice for large-scale graphs where more complex algorithms might be computationally prohibitive. This characteristic is particularly advantageous in real-world applications where quick, near-optimal solutions are preferred over exact, but computationally expensive, methods.

However, the performance guarantee of First Fit in sum coloring is not as robust as in traditional graph coloring. While First Fit is known to use at most  $\Delta+1$  colors for any graph (where  $\Delta$  is the maximum degree of the graph) in standard coloring, its behavior in sum coloring can be significantly suboptimal. Our analysis indicates that the sum of colors assigned by First Fit can be substantially higher than the optimal

sum. This discrepancy arises because the algorithm does not take into account the cumulative effect of previously assigned colors, leading to potentially large sums in densely connected regions of the graph.

Furthermore, the performance of First Fit is highly dependent on the order of vertex presentation. In adversarial or poorly chosen orderings, the sum of the colors can be much larger than the optimal.

For instance, in worst-case scenarios, the sum coloring obtained by First Fit can be exponentially larger than the optimal sum. This sensitivity to vertex ordering highlights a critical limitation of the algorithm, suggesting that improvements or alternative strategies might be necessary for better performance in specific instances.

Despite these limitations, First Fit can still provide valuable approximations in many practical scenarios. Its performance can be significantly improved with preprocessing techniques such as vertex reordering based on heuristics that consider the structure of the graph. For example, ordering vertices by degree or using a breadth-first search ordering can lead to more balanced color assignments and lower total sums. Additionally, hybrid approaches that combine First Fit with other optimization techniques could further enhance its efficacy.

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